

# The Hierarchical Traveling Salesman Problem: Some Worst-Case Results

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# Introduction to the HTSP

- Consider the distribution of relief aid
  - E.g., food, bottled water, blankets, or medical packs
- The goal is to satisfy demand for relief supplies at many locations
  - Try to minimize cost
  - Take the urgency of each location into account

# A Simple Model for Humanitarian Relief Routing

- Suppose we have a single vehicle which has enough capacity to satisfy the needs at all demand locations from a single depot
- Each node (location) has a known demand (for a single product called an aid package) and a known priority
  - Priority indicates urgency
  - Typically, nodes with higher priorities need to be visited before lower priority nodes

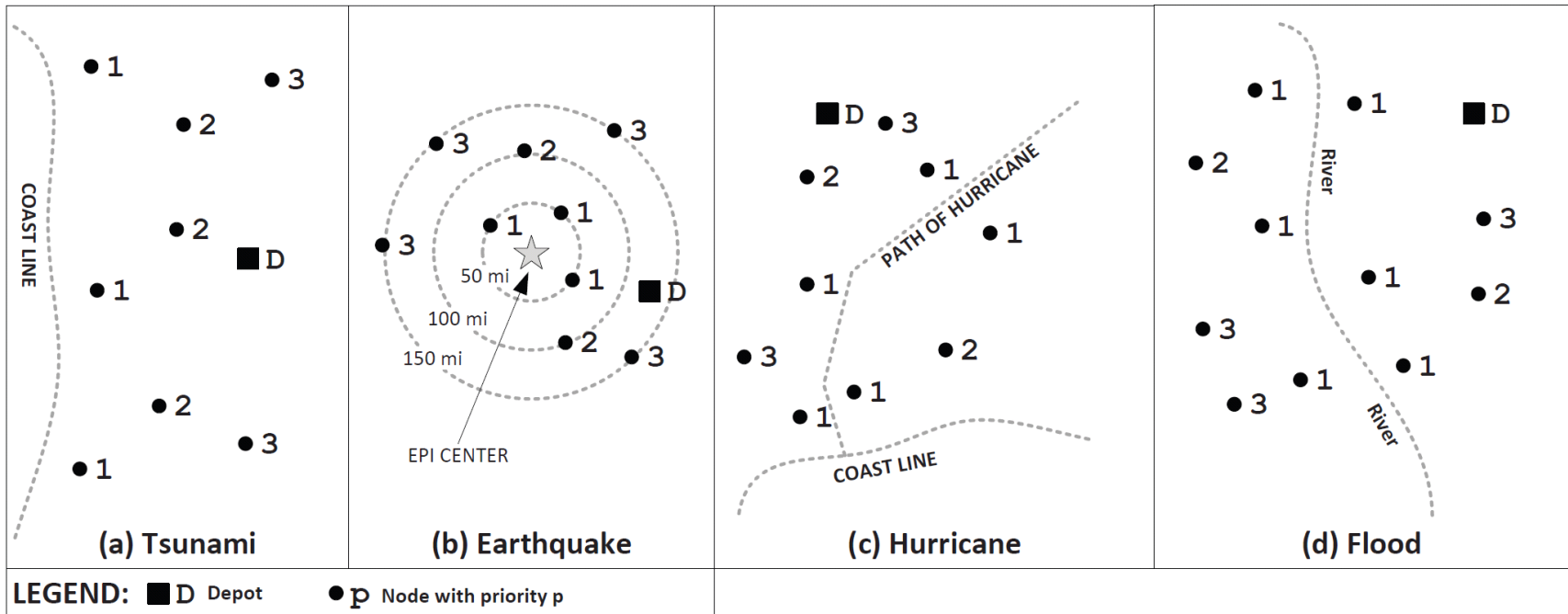
# Node Priorities

- Priority 1 nodes are in most urgent need of service
- To begin, we assume
  - Priority 1 nodes must be served before priority 2 nodes
  - Priority 2 nodes must be served before priority 3 nodes, and so on
  - Visits to nodes must strictly obey the node priorities

# The Hierarchical Traveling Salesman Problem

- We call this model the Hierarchical Traveling Salesman Problem (HTSP)
- Despite the model's simplicity, it allows us to explore the fundamental tradeoff between efficiency (distance) and priority (or urgency) in humanitarian relief and related routing problems
- A key result emerges from comparing the HTSP and TSP in terms of worst-case behavior

# Four Scenarios for Node Priorities



# Literature Review

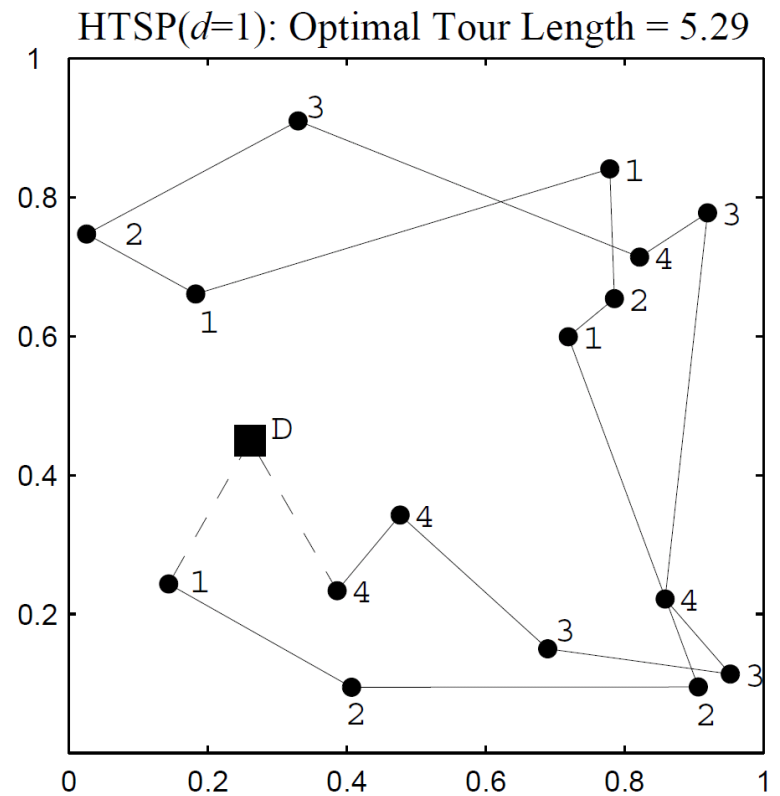
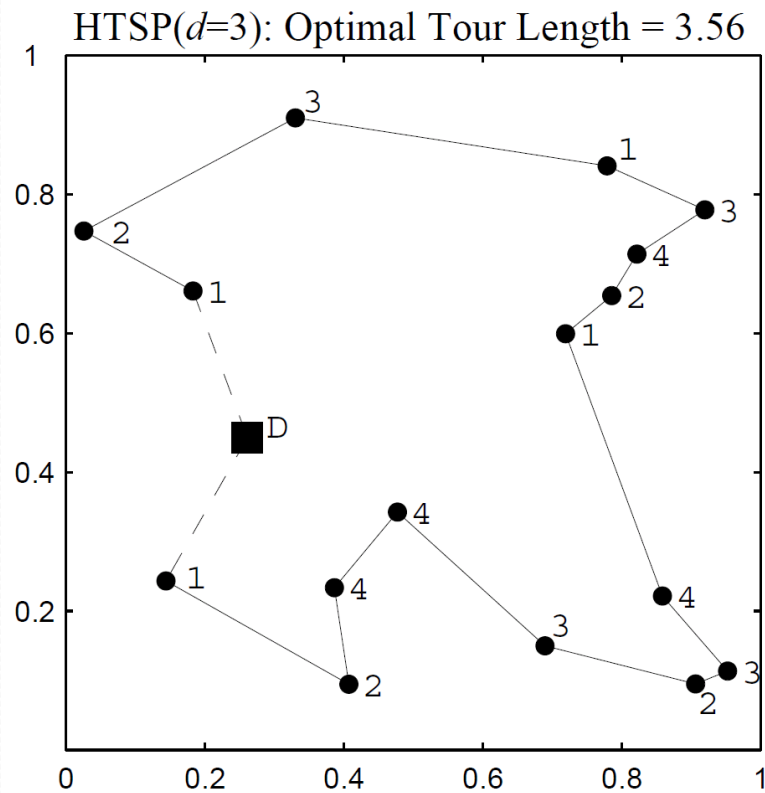
- Psaraftis (1980): precedence constrained TSP
- Fiala Tomlin, Pulleyblank (1992): precedence constrained helicopter routing
- Guttman-Beck et al. (2000): clustered traveling salesman problem
- Campbell et al. (2008): relief routing
- Balcik et al. (2008): last mile distribution

# A Relaxed Version of the HTSP

- Definition: The  $d$ -relaxed priority rule adds operational flexibility by allowing the vehicle to visit nodes of priority  $\pi+1, \dots, \pi+d$  (if these priorities exist in the given instance) but not priority  $\pi+d+\ell$  for  $\ell \geq 1$  before visiting all nodes of priority  $\pi$  (for  $\pi=1, 2, \dots, P$ )
- When  $d=0$ , we have the strict HTSP
- When  $d=P-1$ , we have the TSP (i.e., we can ignore node priorities)



# Efficiency vs. Priority



# Main Results *(Optimization Letters, forth.)*

- Let  $P$  be the number of priority classes
- Assume the triangle inequality holds
- Let  $Z_{d,P}^*$  and  $Z_{TSP}^*$  be the optimal tour length (distance) for the HTSP with the  $d$ -relaxed priority rule and for the TSP (without priorities), respectively
- We obtain the following results below (and the bounds are tight)

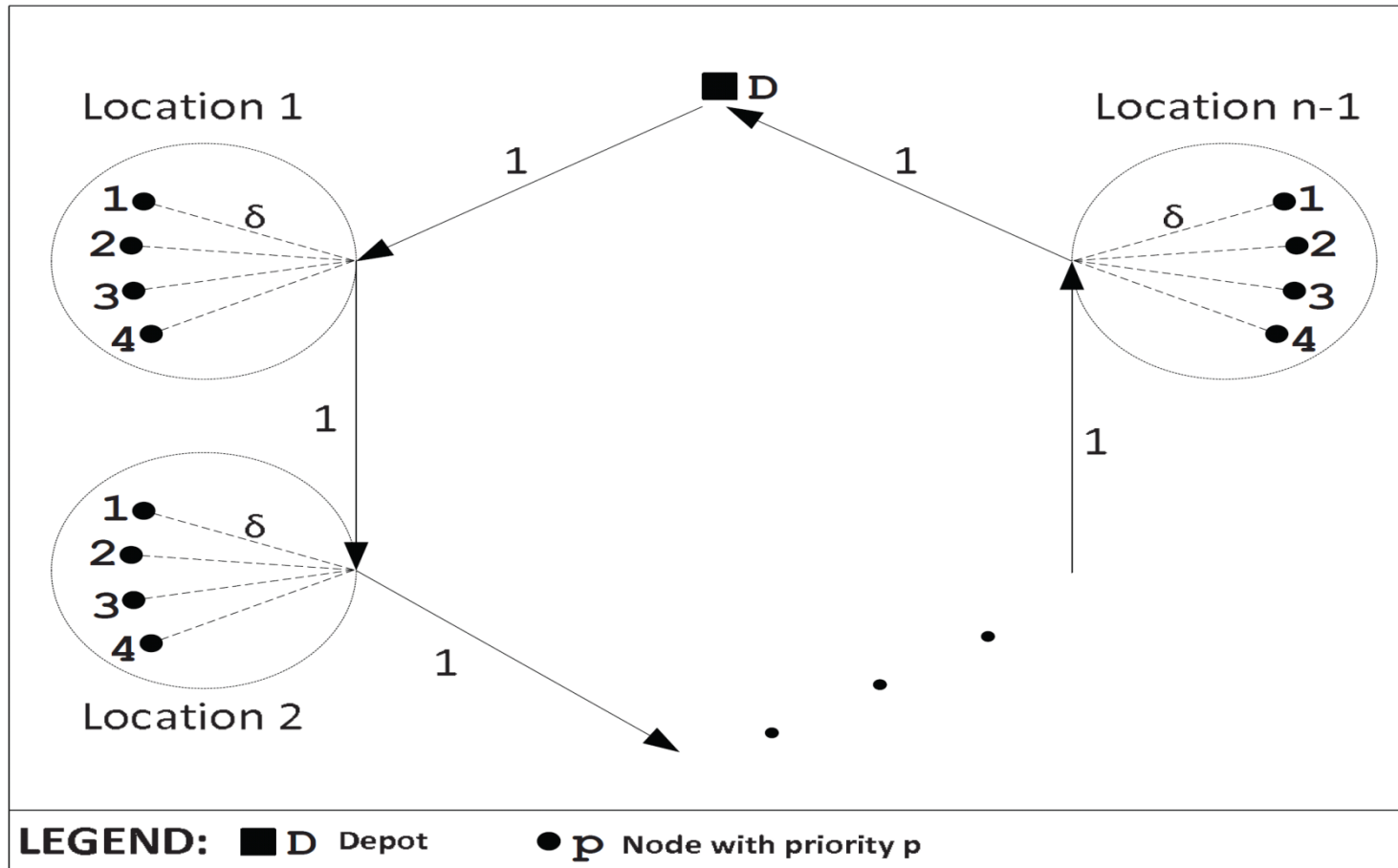
$$(a) \quad Z_{0,P}^* \leq P Z_{TSP}^*$$

$$(b) \quad Z_{d,P}^* \leq \left\lceil \frac{P}{d+1} \right\rceil Z_{TSP}^*$$

# The General Result and Two Special Cases

- $Z^*_{d,P} \leq \left\lceil \frac{P}{d+1} \right\rceil Z^*_{\text{TSP}}$
- If  $d=0$ , we have part (a)
- If  $d=P-1$ , then  $Z^*_{d,P} = Z^*_{\text{TSP}}$

# Worst-Case Example



# Several Observations

- Observation 1. The worst-case example shows that the bounds in (a) and (b) are tight and cannot be improved
- Observation 2. We can “solve” a TSP over the entire set of nodes using our favorite TSP heuristic and obtain a feasible tour for the HTSP by traversing the TSP tour back and forth
- Observation 3. Suppose we select Christofides’ heuristic and let  $Z_{d,P}^h$  be the length of the resulting feasible solution to the HTSP, then we have  $Z_{d,P}^h \leq \frac{3}{2} \cdot \left\lceil \frac{P}{d+1} \right\rceil Z_{TSP}^*$

# Observations and Extensions

- Observation 4. The HTSP (with  $d=0$ ) can be modeled and solved as an ATSP
- Observation 5. Other applications of the HTSP include routing of service technicians, routing of unmanned aerial vehicles, and vehicle routing with backhauls
- We can obtain similar worst-case results (with tight bounds) for the HTSP on the line and the Hierarchical Chinese Postman Problem (HCPP)

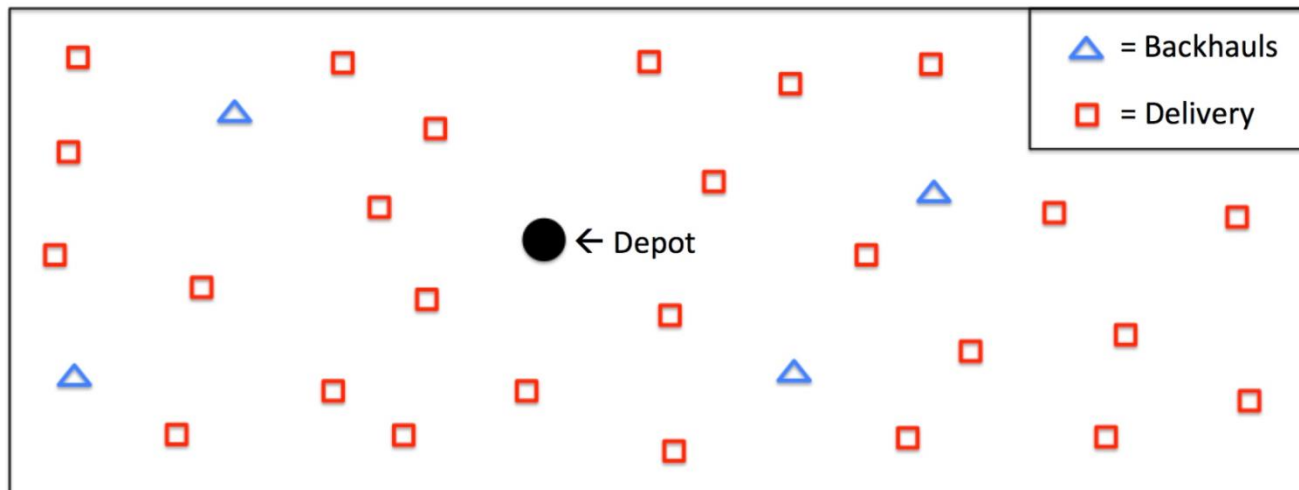
# Vehicle Routing with Backhauls

- Problem statement

- Find a set of vehicle routes that services the delivery and backhaul (pickup) customers such that vehicle capacity is not violated and total distance traveled is minimized
- Such a customer mix occurs in many industries (e.g., the grocery industry)
  - ❖ supermarkets – delivery points
  - ❖ poultry processors, fruit vendors – backhauls

# Vehicle Routing with Backhauls

- Backhauls are serviced at the end of a route
  - Deliveries are high-priority stops
  - Small number of backhauls
  - Difficult to rearrange on-board load in rear-loaded vehicles





# What Next?

- We tried to generalize the previous worst-case bound
- What is the worst-case bound for

$$\frac{Z_{d_1, P}^*}{Z_{d_2, P}^*} \quad \text{where } d_1 < d_2 < P-1?$$

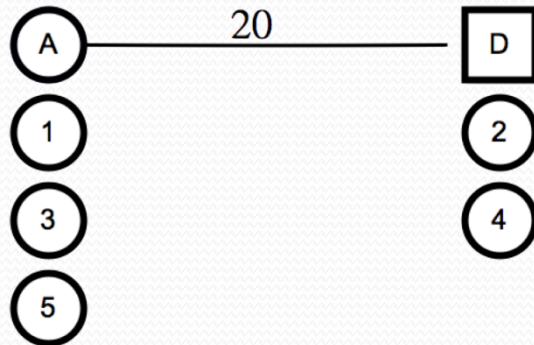
- We have only been able to derive a partial result
- New result (Xiong & Golden, 2013): For any HTSP problem where the triangle inequality holds,

$$\boxed{Z_{0, P}^* \leq 3Z_{1, P}^*}$$

and the bound is tight

# New HTSP Result

- We developed a worst-case example, but a simpler one was found by two of my students (Kim & Park, 2013)
- Consider the example below with  $P=5$



$$Z_{0,5}^* = 120$$

$$Z_{1,5}^* = 40$$

$$Ratio = 3$$

- I would be happy to see someone at this conference extend our work

# Further Remarks

- The HTSP and several generalizations have been formulated as mixed integer programs
- HTSP instances with 30 or so nodes were solved to optimality using CPLEX
- Future work
  - New worst-case results
  - The Hierarchical Vehicle Routing Problem (HVRP)
  - A multi-day planning horizon
  - Uncertainty with respect to node priorities

# Reference

- "The Hierarchical Traveling Salesman Problem" (by K. Panchamgam, Y. Xiong, B. Golden, B. Dussault, and E. Wasil), forthcoming in *Optimization Letters*, 2013.