



keyvb

or not it will halt within the bounds of  $t$  steps (some finite upper bound.) Recall the Unsolvability of Halting Problem (Alternative Proof) handout. One can construct (without contradiction) a semi-tester  $E_a$  that halts iff  $(P,i)$  halts (for any  $(P,i)$ ) For in the 'diagonal' case in which  $E_a$  is fed  $d(E_a)$  (a description of itself) then this states  $E_a$  halts iff  $E_a$  halts (a tautology). The CHP is simply a more restrictive version of semitester  $E_a$  (i.e., falls within  $E_a$ 's extension, or the set of all  $(P,i)$  such that  $(P,i)$  halt.) In other words, let  $S_{Halt} = \{(P,i) \mid (P,i) \text{ halt}\}$ . Let  $S_{Halt}_t = \{(P,i) \mid (P,i) \text{ halt within } t \text{ steps}\}$ . Clearly,  $S_{Halt}_t$  is a subset of  $S_{Halt}$ , so there would exist a modified semitester ( $E_a +$  (a counter) .. call it  $E_{(a,t)}$  whose output comprises the following cases:

case 1: All inputs from the set:  $S_{Halt}_t$ . Then  $E_{(a,t)}$  halts and outputs  $Y$

case 2: All inputs from the set:

$S_{Halt} - S_{Halt}_t = \{(P,i) \mid (P,i) \text{ halt in } (t+1) \text{ steps or more}\}$

Then  $E_{(a,t)}$  halts and outputs  $N$

case 3: Any input from a set  $NoHalt = \{(P,i) \mid (P,i) \text{ don't halt}\}$

Then, (trivially) since any such input would continue past  $t$  steps, then, as in case 2,  $E_{(a,t)}$  halts and outputs  $N$